

# Internal **relational** parametricity, without an interval

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TYPES 2024  
Copenhagen  
11 June 2024

# Parametricity

- ▶  $(A : \text{Type}) \rightarrow A \rightarrow A$
- ▶  $(N : \text{Type}) \rightarrow N \rightarrow (N \rightarrow N) \rightarrow N$
- ▶ Preservation of relations: Reynolds (1983), Wadler (1989), Plotkin–Abadi (1993),  
Bernardy–Jansson–Paterson (2010)
- ▶  $f : (A : \text{Type}) \rightarrow A \rightarrow A$  preserves predicates:

$$f^P : (A : \text{Type})(P : A \rightarrow \text{Type})(a : A) \rightarrow P a \rightarrow P(f A a)$$

fix an  $A$  and an  $a : A$ , then :

$$f^P A (\lambda x. x = a) a \text{ refl} : f A a = a$$

# Internal parametricity

- ▶ Pioneered by Bernardy–Moulin (2012)

- ▶ New syntax:

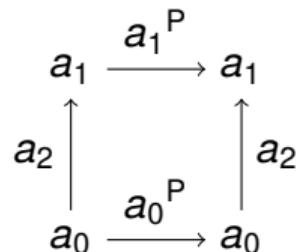
$$\frac{A : \text{Type}}{A^P : A \rightarrow A \rightarrow \text{Type}} \quad \frac{a : A}{a^P : A^P aa}$$

$$(A \rightarrow B)^P f_0 f_1 = (a_0 : A)(a_1 : A) \rightarrow A^P a_0 a_1 \rightarrow B^P (f_0 a_0)(f_1 a_1)$$

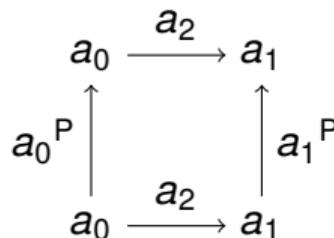
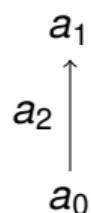
$$\text{Type}^P A_0 A_1 = A_0 \rightarrow A_1 \rightarrow \text{Type}$$

- ▶ Higher dimensional cubes appear when we iterate  $-^P$ :

$$\frac{a_2 : A^P a_0 a_1}{a_2^P : (A^P a_0 a_1)^P a_2 a_2} \\ a_2^P : A^{P^P} a_0 a_0 a_0^P a_1 a_1 a_1^P a_2 a_2$$



$$\frac{(\lambda a.a^P) : (a : A) \rightarrow A^P aa}{(\lambda a.a^P)^P a_0 a_1 a_2 : (A^P aa)^P a_0^P a_1^P} \\ (\lambda a.a^P)^P a_0 a_1 a_2 : A^{P^P} a_0 a_1 a_2 a_0 a_1 a_2 a_0^P a_1^P$$



# Dealing with higher dimensional cubes

paper	substructural interval	model	rel/span
Bernardy–Moulin (2012)	no		rel
Bernardy–Moulin (2013)	yes		rel
Bernardy–Coquand–Moulin (2015) Reboullet (2024)	yes	Reedy fibrant presheaf	rel
Nuyts–Devriese (2018)	yes	ordinary presheaf	rel
Cavallo–Harper (2021)	yes	ordinary presheaf	rel
Altenkirch–Chamoun–Kaposi– Shulman (2024)	no	ordinary presheaf	span

- Relation vs. span

$$A^P : A \rightarrow A \rightarrow \text{Type} \quad \forall A : \text{Type} \quad \text{together with} \quad A \xleftarrow{0_A} \forall A \xrightarrow{1_A} A$$

- In this talk we define a relation-based variant of the ACKS theory.

## An excerpt of the rules

$$\frac{B : \text{Type}}{\forall B : \text{Type} \\ k_B : \forall B \rightarrow B}$$

$$\frac{A : B \rightarrow \text{Type} \\ b_2 : \forall B}{\text{Br}_A b_2 : A(0_B b_2) \rightarrow A(1_B b_2) \rightarrow \text{Type}}$$

- ▶ like Bernardy–Jansson–Paterson (2010)
- ▶ like cubical type theory: “ $\forall B = \mathbb{I} \rightarrow B$ ”

$$\frac{f : A \rightarrow B}{\text{ap } f : \forall A \rightarrow \forall B} \quad \frac{a : (b_x : B) \rightarrow A b_x}{\text{apd } a : (b_2 : \forall B) \rightarrow \text{Br}_A b_2(a(0_B b_2))(a(1_B b_2))}$$

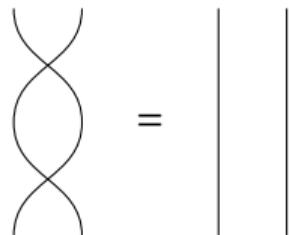
$$\frac{A : \text{Type}}{\text{R}_A : A \rightarrow \forall A} \quad \frac{A : \text{Type}}{\text{S}_A : \forall(\forall A) \rightarrow \forall(\forall A)}$$

# Some equations on $k$ , $R$ , $S$

$$\overline{k_A(R_A a)} = a$$

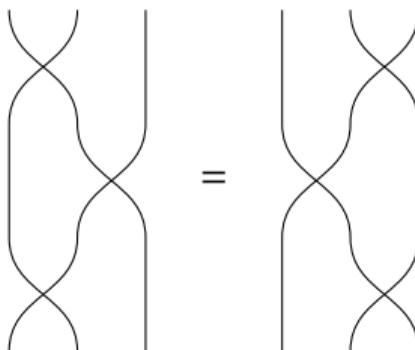
$$\overline{S_A(R_{\forall A} a_2)} = \text{ap } R_A a_2$$

$$\overline{S_A(S_A a_{22})} = a_{22}$$



A string diagram consisting of two strands. The left strand forms a loop that crosses over the right strand. This crossing pattern repeats twice. To the right of the strands is an equals sign (=), followed by two vertical lines.

$$\overline{S_{\forall A}(\text{ap } S_A(S_{\forall A} a_{222}))} = \text{ap } S_A(S_{\forall A}(\text{ap } S_A a_{222}))$$



A string diagram consisting of two strands. The left strand forms a loop that crosses over the right strand. This crossing pattern repeats twice. To the right of the strands is an equals sign (=), followed by another string diagram where the strands cross twice, but the order of crossings is swapped compared to the first diagram.

# Computing Bridge

$$\overline{\text{Br}_{A \circ f} c_2 = \text{Br}_A (\text{ap } f c_2)}$$

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$$\text{Br}_{\lambda c. (a:A) \rightarrow B (c,a)} c_2 f_0 f_1 \cong (a_0 : A(0 c_2))(a_1 : A(1 c_2))(a_2 : \text{Br}_A c_2 a_0 a_1) \rightarrow \text{Br}_B (c_2, a_2) (f_0 a_0) (f_1 a_1)$$

$$\text{Br}_{\lambda X.X} : \text{Br}_{\lambda \text{-Type}} * A_0 A_1 \leftrightarrow A_0 \rightarrow A_1 \rightarrow \text{Type} : \text{Gel}$$

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$$\text{gel} : R a_0 a_1 \cong \text{Br}_{\lambda X.X} (\text{Gel } R) a_0 a_1 : \text{ungel}$$

## Polymorphic identity example

Assume

$$f : (w : \Sigma(A : \text{Type}).A) \rightarrow \pi_1 w,$$

$$A : \text{Type},$$

$$P : A \rightarrow \text{Type},$$

$$a : A,$$

$$p : P a,$$

using the unary version of our theory, we have

$$\text{ungel}(\text{apd } f (A, \text{Gel } P, a, \text{gel } p)) : f (A, a) = a.$$

when setting

$$P := \lambda x.x = a \quad \text{and} \quad p := \text{refl},$$

# Summary

- ▶ A type theory with internal parametricity:
  - ▶ relational (indexed)
  - ▶ there is no substructural interval
  - ▶ parametricity relations compute up to isomorphism, except for the universe:
    - ▶ up to weak section-retraction
- ▶ We use telescopes for  $\forall$ , see the abstract for details.
- ▶ Any model of our previous span-based theory is a model of the relation-based theory.
- ▶ Implementation: [github.com/mikeshulman/narya](https://github.com/mikeshulman/narya)
- ▶ Future work:
  - ▶ stricter theory
  - ▶ H.O.T.T.: adding transport