

Definitely Not Flawed or Not Definitely Flawed?

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Henrik Nilsson

Joint work with John Capper

School of Computer Science

University of Nottingham

How to Design a Type System? (1)

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- This underpins notions such as progress and preservation; it would be strange if a flawless program turns out to be flawed once it is run:
“Well-typed programs do not go wrong [and should therefore be accepted]”
- Of course, depending, we may have to settle for a quite weak notion of “flawless”.

How to Design a Type System? (2)

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- Assume main design goal is to catch as many problems as possible as early as possible.
- Then what about opting for
“the program is not definitely flawed”
if that allows a **stronger** notion of “flawed”,
meaning **more** problems can be caught?
- Of course, a not definitely flawed program
may be revealed to be flawed; e.g. when run.

How to Design a Type System? (3)

- One may question if we then are entitled to talk about a “type system”; Pierce (2002):
A type system is a tractable syntactic method for proving the absence of certain [undesirable] program behaviors by classifying phrases according to the kinds of values they compute.

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- One may question if we then are entitled to talk about a “type system”; Pierce (2002):
A type system is a tractable syntactic method for proving the absence of certain [undesirable] program behaviors by classifying phrases according to the kinds of values they compute.
- What we are considering here is proving the **presence** of undesirable program behaviours, and rule out the program if so.

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- A kind of mixed static and dynamic approach to typing.
- Possible to view as instance of other “mixed approaches”, such as Hybrid Type Checking (Flanagan) or Gradual Typing (Siek and Taha)?

This Talk

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- Overview of the equation-system aspect of a type system for Functional Hybrid Modelling (FHM) as an example of opting for “not definitely flawed”.
- Possible points for discussion:
 - Pros and cons of the approach.
 - What meta-theoretical properties can or should one prove?
 - Other instances of (aspects of) type systems like this?

The FHM Type System

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 - Progress and preservation holds for the **unrefined** system.
- The refinements impose **additional** constraints related to “well-formedness” of modular systems of equations:
 - The refined system rules out **strictly more** programs than the unrefined one.
 - In particular, definitely ill-formed equation system fragments are rejected.
 - But there is **no guarantee** that the resulting system of equations will be well-formed.

The FHM Setting (1)

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- Undirected equations: **non-causal modelling**. (Differential Algebraic Equations, DAE)
- Two-level design:
 - **equation level** for modelling components
 - **functional level** for spatial and temporal composition of components

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- Spatial composition: signal relation **application**; enables modular, hierarchical, system description.
- Temporal composition: **switching** from one structural configuration into another.

The FHM Setting (3)

resistor :: Resistance \rightarrow SR (Pin, Pin)
resistor $r = \mathbf{sigrel} (p, n)$ **where**
 twoPin $\diamond (p, n, u)$
 $r \cdot p.i = u$

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Parametrised model represented by *function* mapping parameters to a model. Note: first class models!

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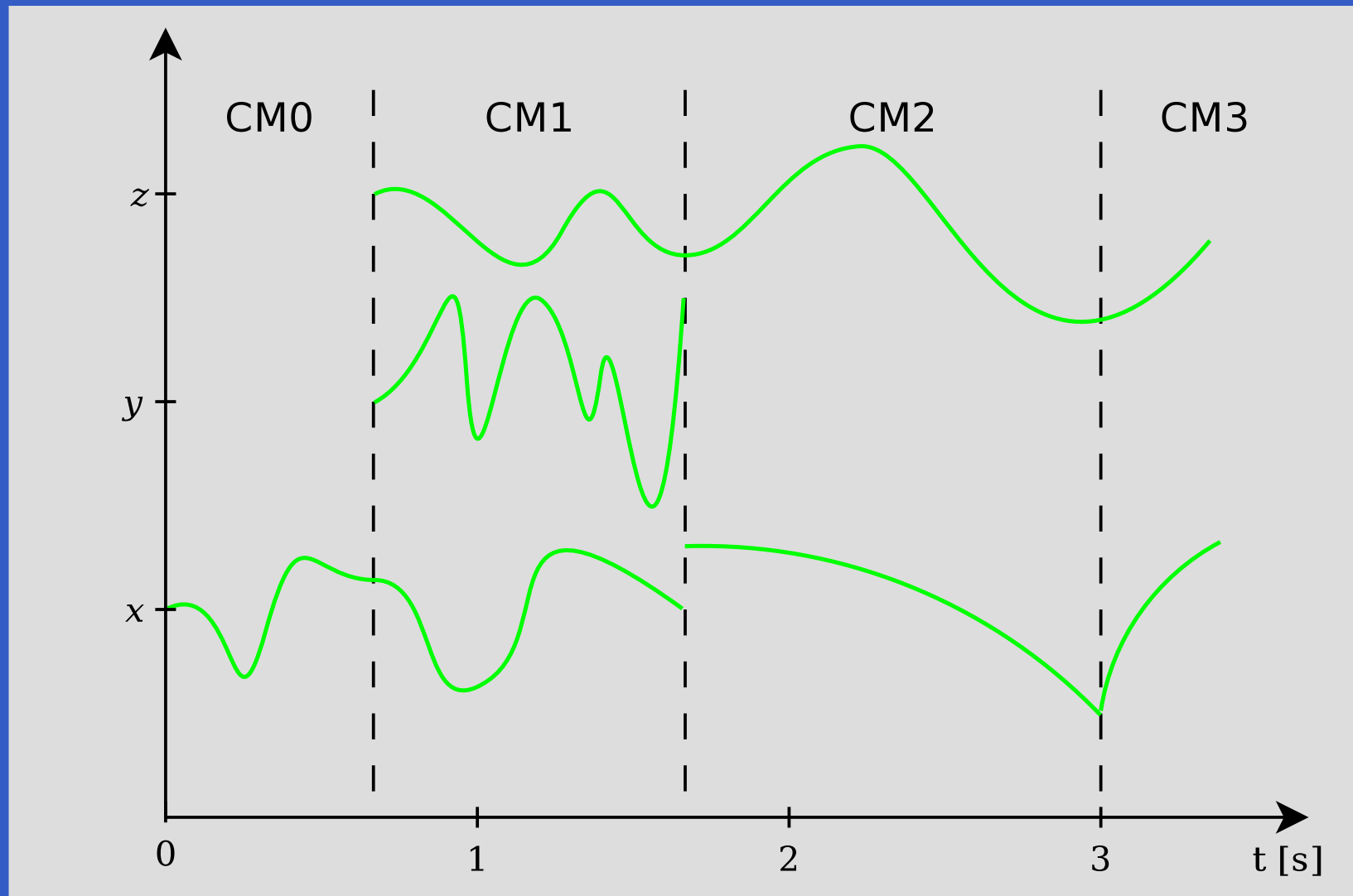
Signal relation **application** allows modular construction of models from component models.

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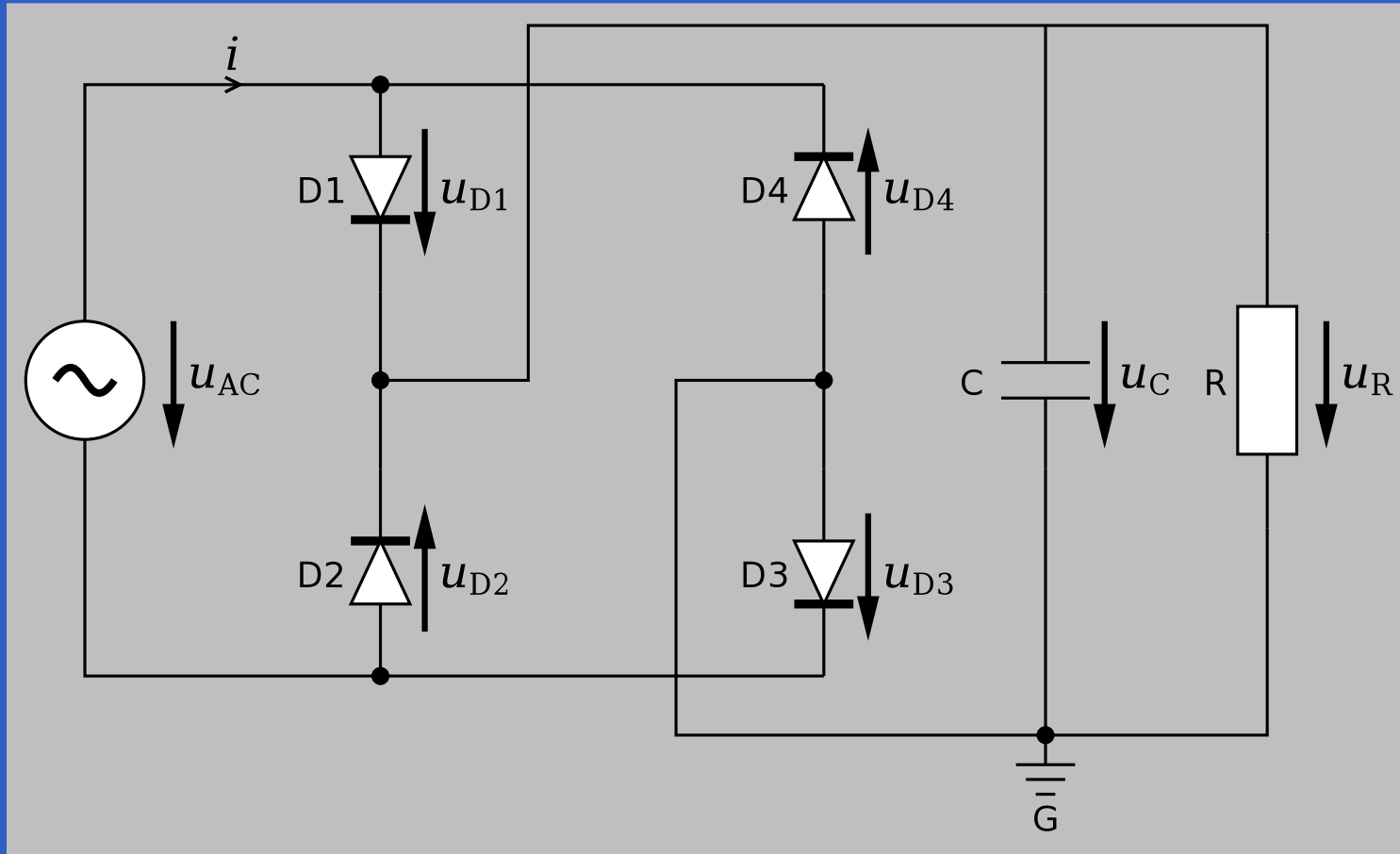
FHM is thus characterised by *iterative staging*:

1. Compute model (“flat” system of equations)
2. Simulate (solve)
3. Repeat

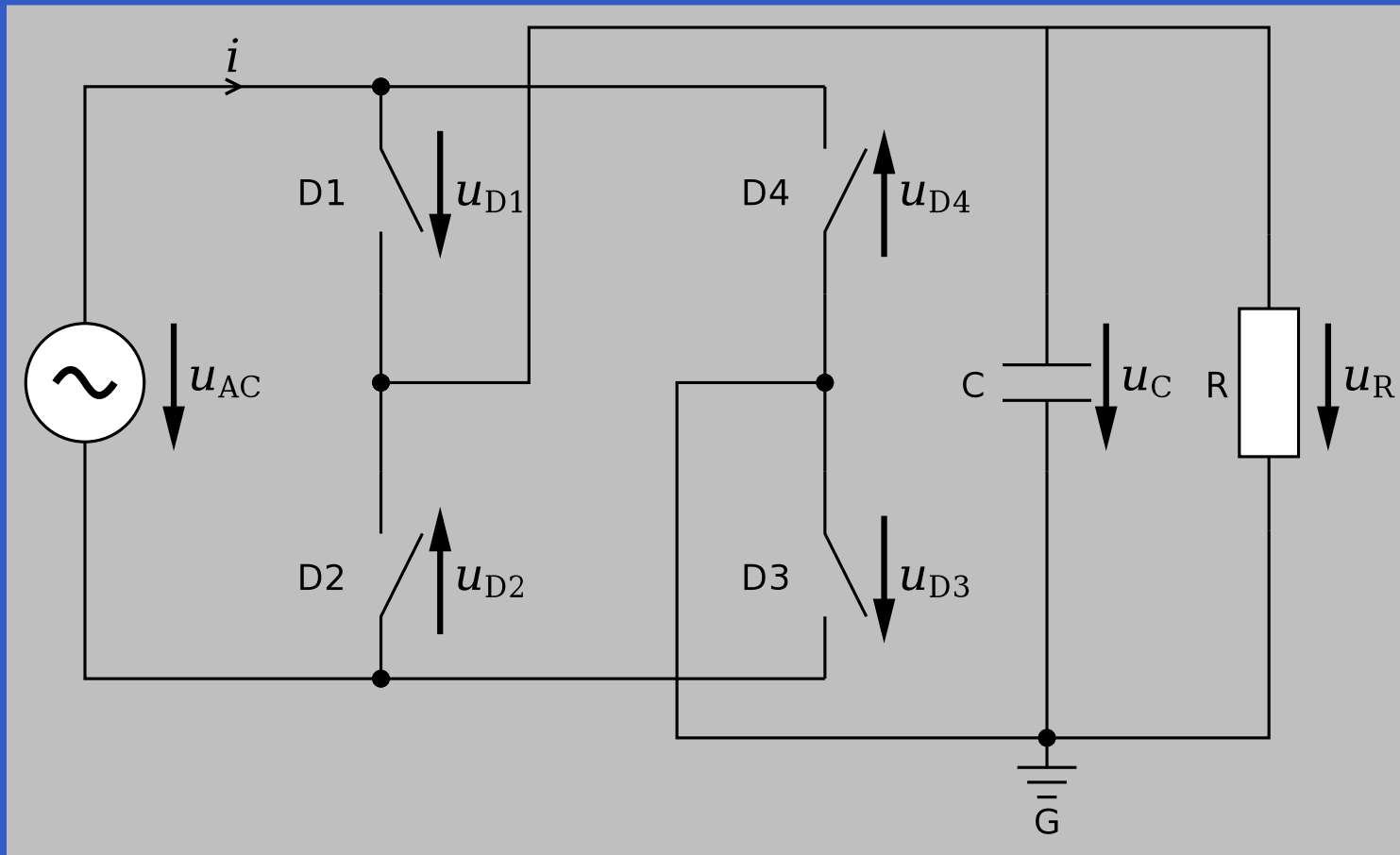
The FHM Setting (5)



Example: Ideal Diodes (1)



Example: Ideal Diodes (2)



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- Can each individual equation system fragment in principle form part of a solvable system?

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- The same holds for:

$$\begin{aligned} x - y + z &= 1 & (2) \\ z &= 2 \end{aligned}$$

The Problem (3)

- Composing (1) and (2):

$$\begin{aligned}x + y + z &= 0 \\x - y + z &= 1 \\z &= 2\end{aligned}\tag{3}$$

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Does have a solution.

- However, the following fragment is over-constrained:

$$\begin{aligned}x &= 1 \\x &= 2\end{aligned}\tag{4}$$

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 - **Structurally non-singular system**
 - **Equation-variable balance**

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- Structural non-singularity is neither a necessary nor sufficient condition for solvability.
- However typical solvers are predicated on the system being structurally non-singular.
- Insisting on structural non-singularity thus makes sense and is not overly onerous.

Structural Non-singularity (2)

Structural singularities can be discovered by studying the *incidence matrix*:

Equations

Incidence Matrix

$$f_1(x, y, z) = 0$$

$$f_2(z) = 0$$

$$f_3(z) = 0$$

$$\begin{matrix} & x & y & z \\ \begin{pmatrix} 1 & 1 & 1 \\ 0 & 0 & 1 \\ 0 & 0 & 1 \end{pmatrix} \end{matrix}$$

Structural Non-singularity (3)

Maybe we can annotate signal relation types with incidence matrices?

$$foo :: SR (Real, Real, Real) \begin{pmatrix} 1 & 1 & 1 \\ 0 & 1 & 1 \end{pmatrix}$$

$foo = \mathbf{sigrel} (x_1, x_2, x_3)$ where

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Yes, possible, but rather complicated and we need to approximate anyway.

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Yes, possible. And preservation holds! (Joey)

But very weak assurances:

$$f(x, y, z) = 0$$

$$g(z) = 0$$

$$h(z) = 0$$

Structural Well-Formedness (1)

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Signal relation types are annotated with **balance**, the number of contributed equations, and types generally carry **constraints** on balance variables:

$$(2 \leq m \leq 4, 3 \leq n \leq 5) \Rightarrow SR (\dots) m \rightarrow SR (\dots) n$$

Structural Well-Formedness (2)

The approach distinguishes between two kinds of variables (number of each kind within parentheses):

- **interface variables** (i_Z)
- **local variables** (l_Z)

and three kinds of equations:

- **interface equations** (i_Q)
- **mixed equations** (m_Q)
- **local equations** (l_Q)

Total number of equations: $a_Q = i_Q + m_Q + l_Q$.

Structural Well-Formedness (3)

A signal relation is **structurally well-formed** (SWF) iff:

1. $l_Q + m_Q \geq l_Z$
2. $l_Q \leq l_Z$
3. $i_Q \leq i_Z$
4. $a_Q - l_Z \leq i_Z$
5. $i_Q \geq 0, m_Q \geq 0, l_Q \geq 0$

The balance (contribution) of a SWF relation is $n = a_Q - l_Z$.

Structural Dynamism

FHM allows for an evolving system of equations by switching blocks of equations in and out:

initially [**when** *condition*₁] \Rightarrow
*equations*₁

when *condition*₂ \Rightarrow
*equations*₂

...

when *condition*_n \Rightarrow
*equations*_n

What about structural well-formedness?

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Exactly one switch-branch active at any point.
How should the number of equations in each
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- **Weak Approach:** same number of equations in each branch.

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Loses too much information.
- **Fair Approach:** branches are reconcilable.

Structural Dynamism and SWF (2)

A switch-block is **reconcilable**, contributing i interface equations, m mixed equations, l local equations, iff i, m, l satisfying the following constraints for each branch k can be found:

$$6. \quad i \geq i_k \geq 0$$

$$7. \quad l \geq l_k \geq 0$$

$$8. \quad m \leq m_k - (i - i_k) - (l - l_k)$$

$$9. \quad i + m + l = i_k + m_k + l_k$$

Note: Interestingly, m may be negative!

Why No Preservation? (1)

Consider:

$foo = \text{sigrel } (x, y)$ where

local z

$$f(x, y, z) = 0$$

$$g(x) = 0$$

$fie = \text{sigrel } (u)$ where

local v

$$foo \diamond (u, v)$$

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$foo \diamond (u, v)$

Both foo and fie are structurally well-formed (why?) with balance 1 and 0, respectively.

Why No Preservation? (2)

But if we carry out some “flattening”:

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Reduction turned a structurally well-formed relation into one that is ill-formed.

Correctness properties?

The following seems plausible:

If $t : C \Rightarrow SR () n$ and $\neg \text{satisfiable}(C)$, then there exists a structural configuration (i.e., a particular choice of switch branches), such that t elaborates to a structurally singular system of equations.

We have not yet attempted to prove this.